Coloring Geographical Threshold Graphs

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Abstract

We propose a coloring algorithm for sparse random graphs generated by the geographical threshold graph (GTG) model, a generalization of random geometric graphs (RGG). In a GTG, nodes are distributed in a Euclidean space, and edges are assigned according to a threshold function involving the distance between nodes as well as randomly chosen node weights. The motivation for analyzing this model is that many real networks (e.g., wireless networks, the Internet, etc.) need to be studied by using a "richer" stochastic model (which in this case includes both a distance between nodes and weights on the nodes). Here, we analyze the GTG coloring algorithm together with the graph's clique number, showing formally that in spite of the differences in structure between GTG and RGG, the asymptotic behavior of the chromatic number is identical: $\chi = \frac{\ln n}{\ln \ln n}(1 + o(1))$. Finally, we consider the leading corrections to this expression, again using the coloring algorithm and clique number to provide bounds on the chromatic number. We show that the gap between the lower and upper bound is within $C \ln n/(\ln \ln n)^2$, and specify the constant C.

1 Introduction

Numerous approaches have been proposed in recent years to study the structure of large real-world technological and social networks, and to optimize processes on these networks. A particularly fertile approach has been to consider the network as an instance of an ensemble, arising from a suitable random generative model. One straightforward example is the random geometric graphs (RGG) model, where nodes are placed uniformly at random in a Euclidean space and edges are placed between any two nodes within a threshold distance. This has the advantage of describing many aspects of systems such as sensor networks, while avoiding unnecessary detail. Even though geometric correlations in RGGs complicate the probabilistic analysis of the model, recent work has clarified many of its structural properties including threshold behavior [10, 5, 6], random walk behavior [1] and chromatic number [8, 9, 10].

RGGs fail, however, to capture heterogeneity in the network. Geographical threshold graphs (GTG) aim at generalizing RGGs, providing this heterogeneity via a richer stochastic model that nevertheless preserves much of the simplicity of the RGG model. GTGs assign to nodes both a location and a weight, which may represent a quantity such as transmission power in a wireless network or influence in a social network. Edges are placed between two nodes if a symmetric function of their weights and the distance between them exceeds a certain threshold [4].

Recent work has analyzed structural properties of GTGs, such as connectivity, clustering coefficient, degree distribution, diameter, existence and absence of the giant component [3, 2]. These properties are not merely of theoretical importance, but also play an important role in applications. In communication networks, connectivity implies the ability to reach all parts of the network. In packet routing, diameter gives the minimal number of hops needed for transmission between two arbitrary nodes. And in the case of epidemics, the existence or absence of the giant component controls whether the epidemic spreads or is contained.

When considering wireless networks, a natural quantity to study is the chromatic number. This is the minimum number of colors needed to color vertices, such that no two adjacent vertices in the graph receive the same color. Treating the colors as the different radio channels or frequencies, the chromatic number gives the minimal number of channels needed so that neighboring radios do not interfere with each other. In this paper we study the asymptotic behavior of the chromatic number for GTGs with constant mean degree. We propose a greedy coloring algorithm, and analyze the behavior of this algorithm along with the graph's clique number. This leads to lower and upper bounds on the chromatic number.

The paper is organized as follows. Section 2 defines the GTG model. Section 3 presents our main asymptotic result, based on our analysis of the coloring algorithm. We show that for graphs G of con-

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stant mean degree, both the clique number $\omega(G)$ and chromatic number $\chi(G)$ are with high probability given by $\frac{\ln n}{\ln\ln n}(1+o(1))$. Section 4 analyzes the gap between lower and upper bounds on the chromatic number, given respectively by the clique number and the greedy coloring algorithm. We show that this gap is within $C\ln n/(\ln\ln n)^2$, and specify the constant C. Finally, Section 5 concludes with open questions regarding the chromatic number for sparser and denser GTGs.

2 Geographical Threshold Graph Model

Given random points $X_1, X_2, \ldots \in [0,1]^2$, chosen iid. uniformly at random, and iid. nonnegative weights $W_1, W_2, \ldots$, we construct the random geographical threshold graphs G_n as follows. Let $N \stackrel{d}{=} \mathrm{Po}(n)$ be the number of the nodes, independent of the X_i and W_i. Let θ_n be a given threshold parameter that depends on the size of the graph n. Then, G_n has vertex set $V(G_n) = \{1, \ldots, N\}$, and an edge $ij \in E(G_n)$ iff

$$\frac{W_i + W_j}{\|X_i - X_j\|^2} \geq \theta_n. \tag{2.1}$$

For technical convenience we identify opposite edges of $[0,1]^2$, making it into a torus.

We will specifically analyze the regime of constant expected degree. If $\mathbf{E}(W_i)$ is a constant, then this occurs when the threshold parameter is linear in the expected number of nodes, $\theta_n = \Theta(n)$. For simplicity we take $\theta_n = n$, since if $\theta_n = cn$ for some constant $c > 0$, the weights can always be rescaled to $W_i := W_i/c$.

3 Asymptotic Results

If G is a graph then $\omega(G)$ denotes its clique number and $\chi(G)$ its chromatic number. We will show formally that the clique number and chromatic number of the geographical threshold graph are essentially the same as those for a random geometric graph with constant average degree:

THEOREM 3.1. *Suppose that $\boldsymbol{Pr}(W_1 > x) = O(x^{-\alpha})$ for some $\alpha > 1$. Then*

$$\frac{\omega(G_n)}{\ln n/\ln\ln n} \to 1 \quad a.s.,$$

and

$$\frac{\chi(G_n)}{\ln n/\ln\ln n} \to 1 \quad a.s.,$$

as $n \to \infty$.

The rest of this section is devoted to proving the theorem.

3.1 Lower bound. Let $\hat{w} \in \mathbf{R}$ be such that $\mathbf{Pr}(W_1 > \hat{w}) \geq \frac{1}{2}$. Then the probability that G_n contains less than $\frac{n}{3}$ vertices with weight more than $\hat{w}$ is exponentially small. Let G'_n be the subgraph of G_n induced by $\frac{n}{3}$ of the points with weights $\hat{w}$ at least. Note that if $i, j \in V(G'_n)$ and $\|X_i - X_j\|^2 < 2\hat{w}/n$ then certainly $ij \in E(G'_n)$. Thus G'_n (and hence also G_n) contains the ordinary random geometric graph $G(\frac{n}{3}, \sqrt{2\hat{w}/n})$ as a subgraph. It follows from computations done in [7] that

$$\mathbf{Pr}(\omega(G_n) < (1-\varepsilon)\ln n/\ln\ln n) = \exp\big(-\Omega(n)\big).$$

3.2 Upper bound. Let us define a "level" L_k as follows:

$$L_{-1} := \{i \leq N : W_i < 1\},$$
$$L_k := \{i \leq N : 4^k \leq W_i < 4^{k+1}\} \quad \text{for } k \geq 0.$$

Note that the set $\mathcal{X}_k := \{X_i : i \in L_k\}$ of the points of the Poisson process corresponding to level k is in fact a Poisson process itself with intensity $n\cdot(F(4^{k+1}) - F(4^k))$ (here F denotes the cdf of W_1) on the unit square and intensity 0 elsewhere. Moreover, these Poisson processes $(\mathcal{X}_k)_k$ are independent.

For $x \in \mathbf{R}^2$ let us denote

$$M_x := \sum_{k=-1}^{\infty} |\{i \in L_k : \|X_i - x\| < 100 \cdot 2^{k+1}/\sqrt{n}\}|,$$

and let us set

$$M := \max_{x \in \mathbf{R}^2} M_x.$$

Then we have the following:

LEMMA 3.1. *The chromatic number satisfies $\chi(G_n) \leq M$.*

Proof. Let us order the vertices by nondecreasing weight and greedily color them. That is, we first color the vertex with smallest weight, then the vertex with second smallest weight and so on; and when we choose a color for a vertex we always pick the smallest possible color (ie. the smallest color that does not occur among the neighbours of the vertex that have already been colored). We claim that in this way we will never need more than M colors.

For ease of notation let us assume (wlog.) that $W_1 \leq W_2 \leq \ldots \leq W_N$. Let us define:

$$N_<(i) := \{j < i : ij \in E(G_n)\}.$$

Note that if $i \in L_k$ and $j \in N_<(i)$ then $\|X_i - X_j\| < 2^{k+2}/\sqrt{n}$. For $1 \leq i \leq N$ let $c(i)$ denote the color that the algorithm has assigned to vertex i.

Now let i be an arbitrary vertex. Let us put $j_0 = i$ and let k_0 denote the level of i. For each of the colors $1, \dots, c(i) - 1$ there is a $j \in N_<(i)$ with $c(j)$ equal to that color. Let $c_1 < c(i)$ be the largest color for which there is no $j \in L_{k_0} \cap N_<(i)$ with $c(j) = c_1$. It is possible that no such c_1 exists, in which case

$$\begin{aligned} c(i) &\le |L_{k_0} \cap N_<(i)| + 1 \\ &\le |\{j \in L_{k_0} : \|X_j - X_i\| < 100 \cdot 2^k/\sqrt{n}\}| \\ &\le M_{X_i} \\ &\le M. \end{aligned}$$

If c_1 exists, then let us pick a $j_1 \in N_<(i) \setminus L_{k_0}$ with $c(j_1) = c_1$. Let k_1 denote the level of j_1. All colors $1, \dots, c_1 - 1$ must occur in $N_<(j_1)$. Let $c_2 < c_1$ be the largest color for which there is no $j \in L_{k_1} \cap N_<(j_1)$ with $c(j) = c_2$. It is possible that no such c_2 exists, in which case

$$\begin{aligned} c(i) &\le |L_{k_0} \cap N_<(i)| + |L_{k_1} \cap N_<(j_1)| + 2 \\ &\le |\{j \in L_{k_0} : \|X_j - X_{j_1}\| < 100 \cdot 2^{k_0}/\sqrt{n}\}| \\ &+ |\{j \in L_{k_1} : \|X_i - X_{j_1}\| < 100 \cdot 2^{k_1}/\sqrt{n}\}| \\ &\le M_{X_{j_1}} \\ &\le M. \end{aligned}$$

Here the first line follows from the fact that each color $\le c(i)$ must either occur as the color of i or j_1 or of a neighbour of i of level k_0 or as the color of a neighbour of j_1 of level k_1; and for the second line we have used the triangular inequality, and the fact that $\|X_i - X_{j_1}\| < 2^{k_0+2}/\sqrt{n}$ and that $\|X_j - X_{j_1}\| < 2^{k_1+2}/\sqrt{n}$ if $j \in N_<(j_1) \cap L_{k_1}$.

Now suppose that $j_1 > \dots > j_m$ and $k_1 > \dots > k_m$ have been defined in such a way that, for $p = 0, \dots, m$, we have $j_{p+1} \in L_{k_p} \cap N_<(j_p)$ and $c(j_{p+1}) < c(j_p)$ is the largest color that does not occur in $\{c(j) : j \in N_<(j_p) \cap L_{k_p}\}$. Let c_{m+1} be the largest color such that there is no $j \in N_<(j_m) \cap L_{k_m}$ with $c(j) = c_{m+1}$. If no such c_{m+1} exists, then

$$\begin{aligned} c(i) &\le \textstyle\sum_{p=0}^m |L_{k_p} \cap N_<(j_p)| + m + 1 \\ &\le \textstyle\sum_{p=0}^m |\{j \in L_{k_p} : \|X_j - X_{j_m}\| < 100 \cdot 2^{k_p}/\sqrt{n}\}| \\ &\le M_{X_{j_m}} \\ &\le M. \end{aligned}$$

The first line follows because necessarily $\{1, \dots, c_m - 1\} \subseteq \{c(j) : j \in N_<(j_m) \cap L_{k_m}\}$ and $c(j_m) = c_m$, $\{c_m + 1, \dots, c_{m-1} - 1\} \subseteq \{c(j) : j \in N_<(j_{m-1}) \cap L_{k_{m-1}}\}$ and $c(j_{m-1}) = c_{m-1}$, and so on. The second line follows because, by the triangle inequality:

$$\begin{aligned} \|X_{j_p} - X_{j_m}\| &\le \sum_{q=p}^{m-1} \|X_{j_q} - X_{j_{q+1}}\| \\ &\le \frac{2}{\sqrt{n}} \sum_{q=p}^{m-1} 2^{k_q+1} \\ &< 2^{k_p+3}/\sqrt{n}, \end{aligned}$$

for all $1 \le p \le m$. And hence, for any $j \in N_<(j_p) \cap L_{k_p}$, we have

$$\begin{aligned} \|X_j - X_{j_m}\| &\le \|X_{j_p} - X_{j_m}\| + \|X_j - X_{j_p}\| \\ &\le (2^{k_p+3} + 2^{k_p+2})/\sqrt{n} \\ &< 100 \cdot 2^{k_p}/\sqrt{n}. \end{aligned}$$

If c_{m+1} exists then we can choose $j_{m+1} \in N_<(j_m) \setminus L_{k_m}$ such that $c(j_{m+1}) = c_{m+1}$ and set k_{m+1} equal to the level of j_{m+1}, and continue by attempting to pick a c_{m+2}. It is clear that the process of picking new c_m's cannot continue indefinitely (certainly there can be no more than N steps), so we can conclude that $c(i) \le M$. Since the vertex i was arbitrary, the claim follows.

To finish the proof of the theorem it now suffices to prove the following lemma:

LEMMA 3.2. *If $\boldsymbol{Pr}(W_1 > x) = O(x^{-\alpha})$ for some $\alpha > 1$ then*

$$\limsup_{n \to \infty} \frac{M}{\ln n / \ln \ln n} \le 1 \qquad a.s.$$

Proof. Let us set

$$M'_x := \sum_{k=-1}^{\infty} |\{i \in L_k : \|X_i - x\| < 200 \cdot 2^{k+1}/\sqrt{n}\}|,$$

and note that if $A := \{(\frac{a}{\sqrt{n}}, \frac{b}{\sqrt{n}}) : 0 \le a, b \le \sqrt{n}\}$ where a and b are integers, then

$$M \le \max_{x \in A} M'_x. \tag{3.2}$$

Let $x \in \mathbf{R}^2$ be arbitrary and note that $M'_x \stackrel{d}{=} \sum_{k=-1}^{\infty} Z_k$, where the Z_k are independent Poisson random variables, and $\mathbf{E}(Z_k) \le \pi(200)^2 \cdot 4^{k+1} \cdot \mathbf{Pr}(W_1 \ge 4^k) = O(4^{k(1-\alpha)})$. So in particular M'_x is itself Poisson with a mean that is bounded above by some constant, μ say. Using a well known bound (see for instance [7]) we see that

$$\begin{aligned} &\mathbf{Pr}(M'_x > (1+\varepsilon)\ln n/\ln\ln n) \\ &\le \left(\frac{e\mu}{(1+\varepsilon)\ln n/\ln\ln n}\right)^{(1+\varepsilon)\ln n/\ln\ln n} \\ &= \exp\big(-(1+\varepsilon+o(1))\ln n\big). \end{aligned} \tag{3.3}$$

Hence, by Eq. (3.2), (by applying the Union bound)

$$\begin{aligned} \mathbf{Pr}\Big(M > (1+\varepsilon)\ln n/\ln\ln n\Big) &\le n e^{-(1+\varepsilon+o(1))\ln n} \\ &\le n^{-\frac{\varepsilon}{2}}. \end{aligned} \tag{3.4}$$

This shows that $M/(\ln n/\ln\ln n)$ is upper bounded by $1+\varepsilon$, whp. To prove an almost sure convergence result, it is possible to adapt a "subsequence trick" from [10], page 123.

4 Mind the Gap

In this section we analyze the gap between lower and upper bounds on the chromatic number, given respectively by the clique number (Subsection 4.1) and the greedy coloring algorithm (Subsection 4.2). In the Subsection 4.3 we show that this gap is within $C \ln n/(\ln\ln n)^2$, and specify the constant C.

4.1 Lower Bound. Informally, we divide the space $[0,1]^2$ into a number of disjoint balls. Then, a clique number of nodes within an arbitrarily chosen ball will give us a lower bound on the chromatic number of the entire geographical threshold graph within $[0,1]^2$. Furthermore, the number of the balls, that is, how we tessellate the space $[0,1]^2$, is a parameter that we discuss later. Formally, the argument is the following.

For some threshold weight w_0, let α be defined by $\mathbf{Pr}(W \le w_0) = \alpha$. We will appropriately choose constants w_0 and α later. Let us define a radius $r_0^2 = w_0/(2\theta_n)$. We consider $b = 1/(2r_0)^2$ disjoint balls with radii r_0, and call these balls B_i. For convenience, tile the square $[0,1]^2$ into $b = 1/(2r_0)^2$ sub-squares of the size $2r_0 \times 2r_0$, and within each of the squares inscribe a ball of radius r_0. The number of nodes within B_i is given by Poisson distribution $\mathrm{Po}(nr_0^2\pi)$, while the number of nodes with weights $\ge w_0$ within B_i is given by $\mathrm{Po}((1-\alpha)nr_0^2\pi)$. For convenience we let $\lambda = (1-\alpha)nr_0^2\pi$. Let us note that for $\theta_n = n$ it follows $b = \frac{1}{4r_0^2} = \frac{\theta_n}{2w_0} = \frac{n}{2w_0}$ (this is $\Theta(n)$) and $\lambda = \frac{\pi}{2}(1-\alpha)w_0$ (this is $\Theta(1)$).

Let us now consider only nodes with weights $W \ge w_0$, that belong to the balls B_i. All nodes with weights $W \ge w_0$ within a ball B_i form a clique, since each pair within B_i satisfies the connectivity relation Eq. (2.1). Let k be a positive integer to be specified later. The number of nodes within B_i satisfies

$$\mathbf{Pr}(\mathrm{Po}(\lambda) \ge k) \ge e^{-\lambda}\frac{\lambda^k}{k!}, \tag{4.5}$$

and we denote $p := e^{-\lambda}\lambda^k/k!$. For I_i being an indicator of the event $\{\mathrm{Po}(\lambda) \ge k\}$, we have $\mathbf{Pr}(I_i = 1) \ge p$. Let us define $J = \sum_{i=1}^{b} I_i$. We will shortly choose k and show that for this choice of k, $\mathbf{Pr}(J = 0) \to 0$. First, $J = 0$ iff all I_i are 0. Second, the indicators I_i are mutually independent, since the balls B_i are mutually disjoint. Then it follows, $\mathbf{Pr}(\cap I_i^c) = \mathbf{Pr}(I_i^c)^b \le (1-p)^b = \exp(\ln(1-p)b)$. We already have $b = \Theta(n)$. We will choose k so that $p = \ln n/n$, which implies $\mathbf{Pr}(J > 0) \ge 1 - \exp(\ln(1-p)b) = 1 - \exp(-\Theta(\ln n)) = 1 - n^{-\Theta(1)}$. Thus, we must solve the following equation in k

$$e^{-\lambda}\frac{\lambda^k}{k!}b = \Theta(\ln n). \tag{4.6}$$

By taking the logarithm, Eq. (4.6) is equivalent to

$$-\lambda + k\ln\lambda - \ln k! + \ln b = \ln(\Theta(\ln n)). \tag{4.7}$$

The Stirling's Formula satisfies $k! = \sqrt{2\pi}k^{k+\frac{1}{2}}e^{-k+\frac{\alpha}{12k}}$ for some $\alpha \in (0,1)$, and applying the logarithm on $k!$ it follows $\ln k! = \frac{1}{2}\ln 2\pi k + k(\ln k - 1) + O(1/k)$. Now, Eq.(4.7) is equivalent to

$$\begin{aligned} k(1+\ln\lambda) + \ln n &= (k+\frac{1}{2})\ln k + \lambda + \frac{1}{2}\ln 2\pi \\ &+ \ln(w_0/2) + O(1/k) + \Theta(\ln\ln n). \end{aligned}$$

Calling $\Lambda = 1 + \ln\lambda$ and $\gamma = \lambda + \frac{1}{2}\ln 2\pi + \ln(w_0/2) + O(1/k) + \Theta(\ln\ln n)$ we obtain the new equivalent equation in k

$$k = \frac{\ln n - \frac{1}{2}\ln k - \gamma}{\ln k - \Lambda}. \tag{4.8}$$

Let us introduce the new variables $y = k/e^{\Lambda}$ and $x = (\ln n - \Lambda/2 - \gamma)/e^{\Lambda}$ and the constant $\eta = 1/(2e^{\Lambda})$. Then Eq. (4.8) is equivalent to

$$y = \frac{x}{\ln y} - \eta. \tag{4.9}$$

For given x and η, Eq. (4.9) has the unique solution in y. It can be easily verified that the solution is given by

$$y = \frac{x}{\ln x}\Big(1 + \frac{\ln\ln x}{\ln x}(1+o(1))\Big). \tag{4.10}$$

Let us call $\Delta = \Lambda/2 + \gamma$. Since $\Delta = o(\ln n)$, $\ln x$ can be expressed as

$$\begin{aligned} \ln x &= \ln(e^{-\Lambda}(\ln n - \Delta)) \\ &= -\Lambda + \ln\ln n + \ln(1 - \Delta/\ln n)) \\ &= \ln\ln n - \Lambda - o(1). \end{aligned} \tag{4.11}$$

Since Λ is constant, it follows

$$\begin{aligned} \ln\ln x &= \ln\Big(\ln\ln n - \Lambda - o(1)\Big) \\ &= \ln\ln\ln n + o(1), \end{aligned} \tag{4.12}$$

and furthermore

$$\begin{aligned} \frac{\ln\ln x}{\ln x} &= \frac{\ln\ln\ln n + o(1)}{\ln\ln n - \Lambda + o(1)} \\ &= \frac{\ln\ln\ln n}{\ln\ln n}(1+o(1)). \end{aligned} \tag{4.13}$$

Now, Eq. (4.9) is equivalent to

$$\begin{aligned} \frac{k}{e^{\Lambda}} &= \frac{1}{e^{\Lambda}}\frac{\ln n - \Delta}{\ln y} - \frac{1}{2e^{\Lambda}} \\ k &= \frac{\ln n - \Delta}{\ln y} - \frac{1}{2}. \end{aligned} \tag{4.14}$$

To proceed, let us calculate $\ln y$

$$\begin{aligned}\ln y &= \ln x - \ln\ln x + \ln\Big(1+\frac{\ln\ln x}{\ln x}(1+o(1))\Big)\\ &= \ln x - \ln\ln x + o(1)\\ &= \ln\ln n - \ln\ln\ln n - \Lambda + o(1).\end{aligned}$$

Plugging the last result into Eq. (4.14), the expression for k finally follows

$$\begin{aligned}k &= \frac{\ln n - \Delta}{\ln\ln n - \ln\ln\ln n - \Lambda + o(1)} - \frac{1}{2}\\ &= \frac{\ln n}{\ln\ln n}\Big(1+\frac{\ln\ln\ln n + \Lambda + o(1)}{\ln\ln n}\Big) - \frac{1}{2}\\ &= \frac{\ln n}{\ln\ln n}\Big(1+\frac{\ln\ln\ln n + \Lambda + o(1)}{\ln\ln n}\Big).\end{aligned}$$

Thus, there is a clique of the size at least k within some ball B_i, with probability $\geq 1 - n^{-\Theta(1)}$. Since $k \leq \omega(G_n) \leq \chi(G_n)$, it follows

$$\frac{\ln n}{\ln\ln n}\Big(1+\frac{\ln\ln\ln n + \Lambda + o(1)}{\ln\ln n}\Big) \leq \chi(G_n).$$

4.2 Upper Bound. In this subsection we derive an upper bound on the chromatic number, given by the greedy coloring algorithm in Section 3. Let us consider the inequality (3.3).

$$\begin{aligned}&\mathbf{Pr}\Big(M'_x > (1+\varepsilon)\ln n/\ln\ln n\Big)\\ &\leq \Big(\frac{e\mu}{(1+\varepsilon)\ln n/\ln\ln n}\Big)^{(1+\varepsilon)\ln n/\ln\ln n}\\ &= \exp\Big\{\Big(B - \ln\big((1+\varepsilon)\frac{\ln n}{\ln\ln n}\big)\Big)(1+\varepsilon)\frac{\ln n}{\ln\ln n}\Big\}\\ &= \exp\Big\{\ln n\Big(\frac{B(1+\varepsilon)}{\ln\ln n} -\\ &- \big(\ln(1+\varepsilon)+\ln\ln n - \ln\ln\ln n\big)\frac{(1+\varepsilon)}{\ln\ln n}\Big)\Big\}\\ &= \exp\Big\{\ln n\Big(\frac{B(1+\varepsilon)}{\ln\ln n} - \frac{(1+\varepsilon)\ln(1+\varepsilon)}{\ln n} -\\ &- (1+\varepsilon) + (1+\varepsilon)\frac{\ln\ln\ln n}{\ln\ln n}\Big)\Big\}\\ &= \exp\Big\{\ln n\Big(\frac{B}{\ln\ln n} + \frac{B\varepsilon}{\ln\ln n} - \frac{\varepsilon(1+o(1))}{\ln\ln n} -\\ &- 1 - \varepsilon + \frac{\ln\ln\ln n}{\ln\ln n} + \varepsilon\frac{\ln\ln\ln n}{\ln\ln n}\Big)\Big\},\end{aligned}$$

where $B = \ln(\mu e)$. Let us choose ε to be

$$\varepsilon = \frac{\ln\ln\ln n + s}{\ln\ln n}, \tag{4.15}$$

then it follows that

$$\begin{aligned}&\mathbf{Pr}(M'_x > (1+\varepsilon)\ln n/\ln\ln n) \leq\\ &\leq \exp\Big\{\ln n\Big(-1+\frac{B-s}{\ln\ln n} +\\ &+ \frac{\varepsilon}{\ln\ln n}\Big(\ln\ln\ln n + B - 1 - o(1)\Big)\Big\}\\ &= \exp\Big\{\ln n\Big(-1+\frac{B-s+o(1)}{\ln\ln n}\Big)\Big\}.\end{aligned}$$

Hence, by Eq. (3.4) and by taking $s \geq B + \delta$ it follows that $\mathbf{Pr}(M < (1+\varepsilon)\ln n/\ln\ln n)$ with probability $\geq 1 - e^{-\frac{\ln n}{\ln\ln n}(\delta - o(1))}$. Thus, for any positive δ, with high probability, that is probability $\geq 1 - e^{-\frac{\ln n}{\ln\ln n}(\delta - o(1))}$, the chromatic number satisfies

$$\chi(G_n) \leq \frac{\ln n}{\ln\ln n}\Big(1+\frac{\ln\ln\ln n + B + \delta + o(1)}{\ln\ln n}\Big).$$

4.3 Comparison of Bounds. Let us now optimize the constants $\Lambda = 1 + \ln\lambda$ and $B = \ln(e\mu) = 1 + \ln\mu$ to minimize the gap between lower and upper bounds on $\chi(G_n)$. We define $s_1 = \max\Lambda$ and $s_2 = \min B$. Thus

$$\begin{aligned}s_1 &= 1 + \max\ln\lambda\\ &= 1 + \max(1-\alpha)nr_0^2\pi\\ &= 1 + \max\ln\frac{\pi}{2}\frac{n}{\theta_n}(1-\alpha)w_0\\ &= 1 + \ln\frac{\pi}{2}\frac{n}{\theta_n} + \max\ln(1-F(w_0))w_0\\ &= 1 + \ln\frac{\pi}{2} + \ln\big(\sup_{w_0\geq 0} w_0(1-F(w_0))\big),\end{aligned} \tag{4.16}$$

by using the definition of $\alpha = \mathbf{Pr}(W \leq w_0)$.

On the other hand $s_2 = 1 + \min\ln\mu$. The conditions imposed on the weight distribution in Lemma 3.2, are $\mathbf{Pr}(W > x) = O(x^{-\alpha})$ for some $\alpha > 1$. Thus, $1 - F(4^j) = O(4^{-\alpha j}) \leq D4^{-\alpha j}$, for an absolute constant D, given by $D = \max_j 4^{\alpha j}(1 - F(4^j))$. Now, we obtain an upper bound on μ

$$\begin{aligned}\mu &\leq \sum_{j=-1}^{\infty}\mathbf{E}(Z_j)\\ &\leq \pi(200)^2\sum_{j=-1}^{\infty}4^{j+1}(1-F(4^j))\\ &\leq \pi(200)^2\sum_{j=-1}^{\infty}4^{j+1}D4^{-\alpha j}\\ &= \pi(200)^2 4D\sum_{j=-1}^{\infty}4^{(1-\alpha)j}\\ &= \pi(200)^2 4D\Big(4^{\alpha-1} + \frac{1}{1-(1/4)^{\alpha-1}}\Big).\end{aligned}$$

That is,

$$s_2 \leq 1 + \ln\Big(\pi(200)^2 4D\Big(4^{\alpha-1} + \frac{1}{1-(1/4)^{\alpha-1}}\Big)\Big). \tag{4.17}$$

Now the lower and upper bounds on $\chi(G_n)$, respectively,

$$\frac{\ln n}{\ln \ln n}\Big(1 + \frac{\ln \ln \ln n + s_1}{\ln \ln n}\Big) \leq \chi(G_n)$$

and

$$\chi(G_n) \leq \frac{\ln n}{\ln \ln n}\Big(1 + \frac{\ln \ln \ln n + s_2}{\ln \ln n}\Big),$$

give us the size of the gap

$$C \ln n/(\ln \ln n)^2. \tag{4.18}$$

Finally, the constant C, specified in the abstract, is

$$C = s_2 - s_1,$$

where s_1 and s_2 are given in Eq. (4.16) and Eq. (4.17), respectively.

5 Conclusion

In this work, we have derived the chromatic number and proposed a coloring algorithm on GTG, for the case of $\theta_n = \Theta(n)$, that is, when the mean degree is constant. It naturally arises, that we are interested into the values of the chromatic number for denser and sparser GTGs. A particularly interesting case would be to show χ around the connectivity regime. The connectivity threshold has been derived to be $\theta_n = \Theta(n/\ln n)$, [2]. However, the methods that we have used here rely heavily on techniques that work for random geometric graphs of equivalent degree. It is unclear whether those techniqes would apply near the connectivity threshold, because the limiting connectivity regime in RGG, when the typical vertex degree grows logarithmically, is of special interest and is already 'hard' [10].

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